

# DSGRN Group Meeting Presentation 2

Adam Zheleznyak

DIMACS REU 2020

June 16, 2020

# Computing linear extensions for Boolean lattices with algebraic constraints

I've been going through the recently submitted paper:

“Computing linear extensions for Boolean lattices with algebraic constraints” (Shane Kepley, Konstantin Mischaikow, Lun Zhang)

In the paper, algorithms are given that can solve this problem:

Given a set of linear polynomials  $\{p_0, \dots, p_K\}$ , an evaluation domain  $\Xi$  that is a polytope, and a partial order  $\prec$  such that  $p \prec q$  when  $p(\xi) < q(\xi)$  for all  $\xi \in \Xi$ , then what are all the total orders

$p_{\sigma(1)} \prec_{\sigma} \dots \prec_{\sigma} p_{\sigma(K)}$  that are a *linear extension* to  $\prec$  such that there is some  $\xi \in \Xi$  where  $p_{\sigma(1)}(\xi) < \dots < p_{\sigma(K)}(\xi)$ ?

# Parameter Space Decomposition Problem

Given an interaction function:

**Definition 1.1.** For  $n \in \mathbb{N}$ , an *interaction function* of order  $n$  is a polynomial in  $n$  variables,  $z = (z_1, \dots, z_n)$ , of the form

$$f(z) = \prod_{j=1}^q f_j(z) \quad (4)$$

where each factor has the form

$$f_j(z) = \sum_{i \in I_j} z_i$$

Say that each  $z_i$  can either evaluate to  $\ell_i$  or  $\ell_i + \delta_i$ . This gives  $2^n$  polynomials of the variables  $\ell_1, \dots, \ell_n, \delta_1, \dots, \delta_n > 0$ . Which total orders are *admissible*: there is some set of values for our variables that such that these polynomials respect that order.

We can use the algorithms as mentioned in the previous slide to get a stricter set of candidates of the admissible linear orders when we look at interaction functions. With some other techniques, this set can be made even stricter.

# Consequences to DSGRN

This gives a major computational improvement:

- Before (2016): DSGRN can handle regulatory networks with nodes that have at most three in-edges and three out-edges.
- Now: DSGRN can handle regulatory networks with nodes that have at most five in-edges (and for some cases of six) and any number of out-edges.

My Goal: To see how well these results can be generalized to handle more values i.e. instead of  $z$  evaluating to either  $l$  or  $l + \delta$ , what about  $l$ ,  $l + \delta_1$ , or  $l + \delta_1 + \delta_2$ ? Or for an arbitrary number of  $\delta$ 's?

